Risk-Hedged Traffic Flow Management under Airspace Capacity Uncertainties

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This paper presents a novel solution to the problem of Traffic Flow Management (TFM) under airspace capacity uncertainty arising from weather or environmental effects. The TFM problem is formulated as a Stochastic Linear Program with multiple available routes between origin-destination pairs, with the weather/environmental factors constraining the probable capacities along these routes. The performance index consists of the delays introduced by deterministic and stochastic capacity constraints. Primary impact of the weather/environmental factors is to require the rerouting of aircraft, causing additional delays. These additional delays require the aircraft to carry additional fuel or to incur costs associated with the failure to meet the schedules in hub-spoke operations. Thus, the variance of delays can be used to define the degree of risk in stochastic TFM, and TFM algorithms that assure the variations in the delay below a specified value can be considered as providing a hedge against uncertain weather or environmental factors. The algorithm developed in this paper produces a risk-hedged TFM decision that results in the least delay at a specified level of acceptable variance. This algorithm represents a dramatic departure from the more traditional stochastic TFM algorithms which minimize expected value of delays, without attempting to control their variances. The performance of the present stochastic TFM algorithm is demonstrated on a Use Case representing NAS operations on a regional scale.

I. Introduction

The impact of stochastic airspace capacity on the National Airspace System (NAS) is a well-studied phenomenon. Uncertainty in airspace capacity originates in the unpredictability associated with atmospheric hazard or climate disruption, which refer to any phenomenon that adversely impacts NAS capacity and its ability to perform efficiently. Examples include volcanic ash, which is an infrequent problem, but can make large areas of the airspace unavailable, or weather events, which are more frequent, but a localized problem. By some estimates, the eruption of the Eyjafjallajökull volcano in Iceland resulted in a loss of $200 million a day for the worldwide aviation industry, and totaled $1.7 billion over the duration of the emission of ashes into the atmosphere. Further losses to trade, due to logistical issues, were not accounted for in these studies. A detailed study of the impact of volcanic ash hazard on the performance and delays in the National Airspace System (NAS) is given in Ref. 3. Although the immediate impact of an atmospheric hazard is on the capacity of the enroute airspace, in some cases, airports or terminal area operations can also be affected. For instance, portions of the airspace are inaccessible due to the potential damage that volcanic ash can cause on aircraft components. Similarly, fog is a known problem at San Francisco Airport affecting the Airport Acceptance Rate and Airport Departure Rate (AAR and ADR respectively) at the airport. Adverse weather can also affect capacity indirectly through an increase in controller workload, possibility of airport closure, and congestion on the airport surface.

Efforts have been made to model the disruption in NAS operations due to many atmospheric hazards, and statistical studies of the NAS have been performed with these models (see Refs. 5-7). If trajectories of all flights and the future capacity of the NAS resources are known exactly, then Traffic Flow Management (TFM) performed is deterministic in nature. Deterministic TFM is well understood (see Refs. 8-21), and computationally feasible

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algorithms are known that can manage NAS traffic while minimizing overall system delay. Under many circumstances, deterministic TFM models cannot accommodate demand (traffic) and capacity uncertainty, resulting in increased delays and conflicts in schedule. Such situations are explored in Ref. 22, where a study evaluated the robustness of a deterministic TFM algorithm, viz. the Bertsimas-Stock Patterson (BSP) model$^{23,24}$, in the presence of uncertainty.

Deterministic TFM optimization methods such as the BSP model obtain TFM directives (e.g. ground hold and airborne delays), but once these directives are incorporated in flight’s schedule, the efficiency of operations in the NAS can be ascertained a priori. However, variability in demand and capacity due to extraneous factors can cause deviation from the nominal performance and result in additional delays.

This paper presents a novel stochastic TFM algorithm that can manage traffic in the presence of uncertainties. In this work, uncertainty in capacity is stressed upon, mirroring the prevalent opinion that one of the major challenges in future TFM will be the ability “to make good decisions in the presence of uncertainty in the prediction of weather”$^{25}$. Results can be extended to address demand uncertainty formulations (see Refs. 26-33). However, demand uncertainty is beyond the scope of this paper.

To motivate the relationship between stochastic weather phenomena and the stochastic TFM, consider Figure 1, which shows probabilistic thunderstorm forecasts from the NOAA, in terms of confidence curves. It is evident that to a route which completely avoids any possible interaction with weather, e.g. the 10%, threshold may require an unreasonably long route. More importantly, even within the region of high probability (where ‘high’ is a set threshold, say 70%), weather is essentially probabilistic, and can be represented by a set of weather realizations, with associated probability values.

A schematic diagram showing three weather scenarios is shown in Figure 2. Each scenario has an associated probability $p_i$ and interacts with the nominal route of an aircraft flying between a given origin / destination pair. The weather phenomena are depicted by blue, red, and green curves, and the nominal route is shown in black.

When a weather scenario is realized, the flight will likely have to modify its route in response to the weather. Therefore, a number of reroutes can be calculated, starting from the end of the common trajectory, which an aircraft flies before any knowledge of the specific weather instance is available (although the statistical data may be available). The reroutes in response to weather are shown as blue, red, and green trajectories in Figure 3.

Depending on which weather scenario is ultimately realized, the flight will follow one of the trajectories. Since trajectories are generally of different length, this causes the entry time of the flight into a downstream Sector to be a stochastic quantity, and consequently, the arrival time at the destination airport is also a stochastic quantity whose stochasticity is linked to the probabilistic behavior of weather. The sample flight times for each scenario are shown as the values $t_1$, $t_2$, and $t_3$ in this figure. A strategy which seeks to minimize the expected value of delay without any consideration of the variance is known as a risk-neutral strategy.

It is possible to follow other strategies. For instance, Figure 4 depicts an intermediate strategy where the flight follows the same weather avoidance route for two scenarios (depicted in blue and green), and a different route for the third scenario. Since the route flown in response to the first two scenarios is longer than the route flown in response to the first scenario in the risk-neutral strategy, the expected value of the delay time increases as shown at the top of Figure 4, but the variation in the delay from the mean decreases. An extreme strategy is shown in Figure 5, where the flight follows the red reroute regardless of which weather scenario is realized, because this route avoids all simulated scenarios. However, the flight time using the red route is the largest among the three. Using this risk-averse strategy, the expected delay is the largest; however, the variance is zero.
Figure 2. Weather Scenarios and Flight Routes in Response to Scenario Realization
Figure 3. Risk Neutral Strategy for Rerouting in Response to Weather Hazard
Figure 4. Intermediate Strategy Resulting in Larger Delays and Smaller Variation
Figure 5. Risk-Averse Trajectory with No Variation

The relationship between expected delay, variation in delay, and choice of strategy motivates a multi-objective tradeoff. Multi-objective problems that balance mean throughput against variance are common in financial engineering, where the mean throughput is interpreted as the expected return-on-investment of a portfolio, and the variance is interpreted as the associated risk. A portfolio in the context of financial engineering is a convex combination of resources with known return and risk, where the decision variables are the proportions with which each resource is selected.

Stochastic TFM with an aim to maximize mean throughput or to minimize mean delays are alone not sufficient for probabilistic TFM. Recent works in the literature (see Refs. 35-39) have addressed the issue of minimizing the expected value of delays using either an LP framework or flow-based optimization. Stochastic Linear Programming has also been used by Clare and Richards to obtain TFM solutions with probabilistic capacity constraints, although Ref. 40 allows for capacity violations, and does not address the resulting variability in performance.

The authors believe that the work presented in this paper is the first to address the management of variance in addition to delay minimization TFM. This is an important result, because – as noted earlier – NAS operations are typically designed for expected value of performance whereas inefficiency is caused by a variation in this value.

This paper is organized as follows. Section II introduces the concept of risk in TFM and describes methods to manage risk using stochastic programming. Section III describes the formulation of a stochastic TFM in the LP framework. Section IV contains details of the Use Case which serves as the feasibility demonstration of the risk-management algorithm. Conclusions, summary, results, and directions for future work are presented in Section V.

II. The Concept of Risk in TFM and Risk-Management Using Stochastic Programming

Since risk is a concept that is used extensively in this work, it requires a definition at the outset. Every TFM algorithm that uses optimization algorithm, for example, those listed in Section I, requires the minimization or maximization of a cost function. For instance, the BSP, BLO, and AAEP models typically minimize the ground and airborne delays, number of cancelled flights, and the amount of deviation along a flight segment; or a
combination of any of the above. Depending on the specific nature of the problem, additional components can also be introduced in the cost function. The TFM algorithm produces a set of directives, e.g. the amount of ground delay/stop or airborne holding and route selection. The optimization is subject to constraints on the dynamics of the flight and the available resources such as Sector capacity.

In the presence of capacity uncertainty, the cost function can take different values in response to a specific realization of the capacity constraint. For instance, if there is a probability associated with a downstream Sector for a flight being blocked, the flight may require a longer reroute. This paper defines Risk as the variation in the cost function in response to stochasticity. It needs to be clarified at the outset that Risk in the context of this work is not the probability that a flight will encounter weather. The TFM problem in this work requires all flights to avoid all given realizations of the weather.

A. Using the Spread of a Distribution to Characterize Risk

![Figure 6. Spread of a Probability Distribution as a Measure of Risk](image)

To further develop the concept of risk, Figure 6 presents a depiction of two probability distributions, both of which assign a probability to the value of an objective function. The first distribution has a considerably smaller “spread”, and if the normal distribution is assumed, this translates into a smaller value of the variance. The implication of the smaller variance is that one can assure with 99.7% confidence that the value of the objective function, i.e. system delay, will be within the 3-sigma range of the expected delay. On the contrary, from the distribution on the right, the 99.7% confidence only exists for a significantly larger range around the expected value. A larger range of uncertainty has implications on efficient operation of airlines, especially in a hub-and-spoke configuration. If an airline is unable to guarantee with a high level of confidence that a flight will reach its destination within a certain range of its expected arrival time, then this may impact connecting flights and turn-around times significantly. Furthermore, the likelihood of an extreme event occurring, e.g. an extremely large deviation from the mean delay, is significantly higher when the spread of the distribution is large.

B. Relationship between Stochastic Weather and Capacity Uncertainty

Adverse weather directly affects NAS performance by limiting resources that may be available to the flights. For instance, fog at the San Francisco Airport often limits the aircraft acceptance or departure rates. In this work, the focus is on enroute capacity uncertainty, although the algorithms and analysis can be generalized to include airport capacity uncertainty.

Although Sector capacity determination is an open area of research (see Refs. 44-49), a starting point for nominal Sector capacity assessment is the Sector’s Monitor Alert Parameter (MAP) value. The uncertainty in Sector capacity is most easily derived by an area-based measure, which calculates the area of overlap between the weather phenomenon in a realization, with the Sectors of interest. This is schematically shown in Figure 7, where a weather-affected area, depicted by the polygon, interacts with a Sector of interest, shaded in gray. Depending on what fraction of the Sector intersects with the weather, and given the probability of occurrence of a scenario, a histogram of capacities, ranging from zero to maximum capacity $c_{\text{max}}$, can be constructed. This procedure is independent of the LP formulation or solution and can be performed offline.

The quality of the histogram depends on the number of scenarios available, but only to an extent, since only integers can be used to represent MAP values.
C. Strategies for Risk Management

Several strategies may be employed to manage risk in TFM; some of which are described here. These are depicted in Figure 8. One possible strategy is to design a route that avoids the mean of all realizations of the weather. However, since there is a finite probability for each realization occurring (possibly based on forecast or simulations), the “mean weather” may never be realized, and a reroute will be required when one of the weather realization occurs. This is shown in the top left of Figure 8, where the planned route is shown in black. Consequently, a reroute will be required, shown in blue. However, solving for a reroute once the flight has already departed on its nominal trajectory will require another TFM optimization problem, and the feasibility of the solution, i.e. a new feasible trajectory may not be available.

The second strategy, shown on the top right, is to calculate a reroute for all weather realizations (shown in blue, red, and green), and travel on the mean path (shown in black). However, this strategy does not mitigate the probability with which adverse weather will be encountered, and the flight will be forced to fall back on Strategy 1, i.e. plan a reroute from the point at which weather impacts the trajectory.

The third strategy, shown on the bottom left, is to travel on a route that avoids all simulated weather realizations. This is known as a ‘risk-averse’ strategy, and can cause unnecessary delays especially if the weather realization did not ultimately require a longer reroute. Furthermore, there are no guarantees on the expected delay, although the variation in the delay may be small.

The fourth strategy, shown on the bottom right, divides the trajectory of an aircraft into two segments. The first segment is known as the first stage, and is common to all routes designed to avoid the weather realization. This is
shown in black in the figure. The second segment, known as the second stage or recourse stage, consists of reroutes designed to avoid each weather scenario provided to the flight. These are shown in blue, red, and green. The first stage trajectory is designed to guarantee feasibility of each of the recourse trajectories. Consequently, the flight does not need to solve another TFM problem if adverse weather is encountered. The strategy for a flight, i.e. the choice of a reroute in response to a weather realization, will affect the expected value of the system delay as well as its variation. When the TFM problem is solved by minimizing the expected value of system delay over all scenarios, this is known as a ‘risk-neutral’ strategy, and is generally posed as a two-stage stochastic linear program. This is described in the next section.

D. Stochastic Linear Program to Solve the Two-Stage Recourse Problem

Stochastic Linear Programs (SLP) are generally modeled as two-stage recourse problems. The uncertainty in the problem is modeled by defining a finite number of scenarios \( \{ \omega_1, \omega_2, ..., \omega_k \} \) with associated probabilities of occurrence \( \{ p_1, p_2, ..., p_k \} \). The general form of a two-stage recourse problem is as follows

\[
\max c^T x + \sum_{i=1}^{k} p_i r_i \\
\text{subject to} \\
x \geq 0 \\
T_i x + W_i y_i \leq h_i \\
\text{where } r_i(x) = \max g_i^T y_i \\
\text{subject to } y_i \geq 0
\]  

(1)  

(2)

In the foregoing equations, \( x \) is the first-stage decision variable that does not respond to the uncertain scenario \( \omega \). It is determined before any information regarding the uncertainty data has been obtained. On the other hand, \( y_i \) is the second-stage decision variable that is determined after deriving the observations regarding the uncertainties. Note that a different recourse decision \( y_i \) will be made when the \( i \)th scenario is realized. The goal of the two-stage model is to identify a first-stage solution \( x \) that is well-positioned against all possible manifestations of the uncertain scenarios in the future. An optimal first-stage solution \( x \) will tend to have the property that leaves the second stage in a position to exploit the advantageous scenarios without excessive vulnerability to disadvantageous scenarios.

The constraint matrix \( A \) and the constraint vector \( b \) are the deterministic constraints that are completely independent of the uncertainty. The objective function coefficients or cost for the first stage variables are given by \( c \). The recourse costs \( r_i(x) \) are a function of the first-stage decision \( x \) and the uncertain data associated with the \( i \)th scenario. The objective function coefficients or costs for the recourse variables are given by \( g_i \). The uncertain constraints are defined by the matrices \( T_i \) and \( W_i \) and vector \( h_i \) known as the technology matrix, recourse matrix, and resource vector, respectively. The two-stage recourse problem takes the following form:

\[
\text{max } \begin{bmatrix} c^T & p_1 g_1^T & p_2 g_2^T & \cdots & p_k g_k^T \end{bmatrix} \begin{bmatrix} x \\ y_1 \\ y_2 \\ \vdots \\ y_k \end{bmatrix} \\
\text{subject to } \begin{bmatrix} A \\ T_1 \end{bmatrix} \begin{bmatrix} x \\ y_1 \\ y_2 \\ \vdots \\ y_k \end{bmatrix} + \begin{bmatrix} W_1 \\ T_2 \\ W_2 \\ \vdots \\ T_n \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_k \end{bmatrix} \leq \begin{bmatrix} b \\ h_1 \\ h_2 \\ \vdots \\ h_k \end{bmatrix} \\
x, y_1, y_2, ..., y_n \geq 0
\]  

(3)

E. Exploring the Pareto Frontier in the Mean-Variance Tradeoff Space

The solution to the linear program given by Eq. (3) seeks to maximize the expected value (mean) of the cost over all possible scenarios. However this linear program does not explicitly consider the variance of the recourse costs over all possible scenarios. Modern Portfolio Theory (MPT) states that the variance can be treated as a measure of risk and that there exists a trade-off in the mean-variance space as shown in Figure 9. Since the linear program given by Eq. (3) maximizes mean without considering the variance it gives the risk-neutral solution shown by red in Figure 9.
If the objective function seeks to minimize the variance of the recourse costs, the solution obtained will be the risk-averse solution shown in green. However, including the variance in the objective function will make the programming problem quadratic. The present research develops a linear programming formulation that generates the Pareto front of optimal solutions in the mean-variance space. Once the Pareto front is generated, the solution that maximizes the mean at a given level of acceptable risk can be determined as shown in orange in Figure 9.

F. Generating the Pareto Frontier by Solving the LP with Reduced Variance

The problem of minimizing the variance in the recourse costs $r_1, r_2, \ldots, r_k$ for various scenarios $i = 1 \ldots k$ can be posed as a linear program similar to Eq. (3) with additional constraints that limit the deviation of the recourse cost from the mean recourse cost.

The first step is to solve the risk-neutral stochastic linear program given by Eq. (3). Let $\bar{r}$ denote the mean of the recourse costs over all possible scenarios for the solution to Eq. (3). The mean recourse cost can be calculated from the scenario costs as shown below:

$$\bar{r} = \sum_{i=1}^{k} p_i r_i$$  \hspace{1cm} (4)

The deviation $d_i$ of the recourse cost for the $i$th scenario from the mean recourse cost is given by the following expression:

$$d_i = r_i - \bar{r}$$  \hspace{1cm} (5)

Let $D = \max_i |d_i|$ denote the maximum deviation over all scenarios. The main idea is to iteratively solve the linear program posed by Eq. (3), by adding constraints which limit the deviation of the recourse costs from the mean to be smaller than the maximum deviation at the previous iteration. Thus the additional constraint to be imposed for the $i$th scenario is given by the following:

$$|r_i - \bar{r}| \leq \alpha D, \quad 0 < \alpha < 1$$  \hspace{1cm} (6)

The set of absolute value constraint shown in the foregoing equation can be replaced by the following two sets of linear constraints:

$$r_i - \bar{r} \leq \alpha D$$

$$\bar{r} - r_i \leq \alpha D$$  \hspace{1cm} (7)

This approach requires $2k$ additional constraints in the stochastic linear program. Since the recourse costs and their mean are linear functions of the decision variables, the constraints can be written as follows:

$$g_1^T y_i - (p_1 g_1^T y_1 + \ldots + p_i g_i^T y_i + \ldots + p_k g_k^T y_k) \leq \alpha D$$

$$-g_i^T y_i + (p_1 g_1^T y_1 + \ldots + p_i g_i^T y_i + \ldots + p_k g_k^T y_k) \leq \alpha D$$  \hspace{1cm} (8)
The modified linear program obtained upon inclusion of the \(2k\) additional constraints is given as follows:

\[
\begin{align*}
\max \{ & c^T \begin{bmatrix} y_1 & y_2 & \vdots & y_k \end{bmatrix} \\
& \begin{bmatrix} x_1 & x_2 & \vdots & x_k \end{bmatrix} \\
& \begin{bmatrix} A \end{bmatrix} \\
& \begin{bmatrix} b \end{bmatrix} \\
& \begin{bmatrix} T_1 \end{bmatrix} \\
& \begin{bmatrix} h_1 \end{bmatrix} \\
& \begin{bmatrix} T_2 \end{bmatrix} \\
& \begin{bmatrix} h_2 \end{bmatrix} \\
& \begin{bmatrix} \vdots \end{bmatrix} \\
& \begin{bmatrix} \vdots \end{bmatrix} \\
& \begin{bmatrix} T_k \end{bmatrix} \\
& \begin{bmatrix} h_k \end{bmatrix} \\
& \begin{bmatrix} W_1 \\ W_2 \vdots \end{bmatrix} \\
& \begin{bmatrix} \vdots \end{bmatrix} \\
& \begin{bmatrix} -p_1 g_1^T \\ (1-p_1) g_1^T \end{bmatrix} \\
& \begin{bmatrix} (p_1 - 1) g_1^T \\ -p_1 g_1^T \end{bmatrix} \\
& \begin{bmatrix} -p_2 g_2^T \\ (1 - p_2) g_2^T \end{bmatrix} \\
& \begin{bmatrix} (p_2 - 1) g_2^T \\ -p_2 g_2^T \end{bmatrix} \\
& \begin{bmatrix} -p_k g_k^T \\ (1 - p_k) g_k^T \end{bmatrix} \\
& \begin{bmatrix} (p_k - 1) g_k^T \\ -p_k g_k^T \end{bmatrix} \} \leq 0
\end{align*}
\]

\[(9)\]

\[x, y_1, y_2, \ldots, y_k \geq 0\]

The next section formulates the risk-hedging stochastic LP for TFM.

### III. Deterministic and Stochastic LP Formulation for TFM

A complete description of the Mixed-Integer Linear Program (MILP) formulation for Stochastic Air Traffic Flow Management Rerouting Problem (SATFMRP) is given in Refs. 41 and 42. In this section, some of the key definitions and constraints are presented. The notation in this paper follows that used in Refs. 41 and 42, and the decision variable of interest will be referred to as the BLO variable.

#### A. Decision Variable and Data Sets

The variable of interest in this formulation is \(x_{f,XY}^j\), a binary variable, i.e. 0 or 1. A value of 1 indicates that flight \(f\) (member of set \(F\)), reaches Node Y from Node X, by time interval \(t \in T\) using an arc connecting the two nodes. Nodes X and Y belong to set \(N_f\) that is composed of all nodes on the route(s) of flight \(f\). The sets \(K^d\) and \(K^a\) denote the set of nodes corresponding to departure and arrival airports, respectively. Since an airport in general is both a departure as well as an arrival airport, \(K^d \cap K^a \neq \emptyset\). Let \(k_f^d \in K^d\) and \(k_f^a \in K^a\) denote departure and arrival airport nodes for flight \(f\), respectively. Node \(q(k_f^d)\) denotes the departure airport boundary, and \(p(k_f^a)\) denotes the arrival airport boundary. The distinction between an airport node and its boundary node is to allow for a detailed model for ground holds and runway delays.

The arc \(XY\) is a member of set \(A_f = \{XY | X, Y \in N_f\}\) that is composed of all arcs on the route(s) of flight \(f\). The set \(\Gamma_f^+(X) = \{Y | XY \in A_f\}\) and \(\Gamma_f^-(X) = \{Y | YX \in A_f\}\) are the set of nodes that have arcs from Node X and leading into Node Y respectively, for flight \(f\).

Whereas sets \(N_f\) and \(A_f\) denote all possible nodes and all possible arcs for flight \(f\), the sets \(N_f^* \subset N_f\) and \(A_f^* \subset A_f\) denote the nodes and arcs corresponding the scheduled route of flight \(f\). The variables \(l_{f,XY}\), \(r_f\), and \(d_f\) denote the travel time (number of time periods) for flight \(f\) over arc \(XY\), the scheduled departure time period, and the scheduled arrival time period, respectively. It is noted in Ref. 42 that \(l_{f,k_f^d,a(k_f^d)} = l_{f,k_f^d,a(k_f^a)} = 0, \forall k_f^d \in K^d, k_f^a \in K^a\). In other words, a flight reaches the departure airport boundary immediately after leaving the departure airport node, and a flight reaches the arrival airport node immediately after leaving the arrival airport boundary. It also follows that

\[
r_f = d_f + \sum_{XY \in A_f^*} l_{f,XY}\n\]

\[(10)\]
In other words, the scheduled arrival time of the flight is given by the sum of the departure time and flight times along scheduled route segments. Furthermore, $T_{f,XY}$ and $I_{f,XY}$ denote the maximum and minimum number of time segments for flight $f$ on arc $XY$.

The 0-1 BLO variables can be used to determine quantities of interest for TFM. For instance, the time segment in which the flight $f$ reaches node $n$ is denoted by $T_{f,n}$ and given by the following summation:

$$T_{f,n} = \sum_{x \in T_{f,xy}} \sum_{t \in T_f} t(x_f^{t-1} - x_f^{t-1})$$

(11)

It follows from Eq. (11) that given the time of entry at a node and the time of entry at a preceding node, the number of time intervals required to travel on the arc connecting the nodes can be calculated. Additional quantities such as sector counts (given the arcs belonging to a sector) can also be calculated, as detailed in Refs. 41 and 42.

**B. Constraint Formulation**

The variables are linked with constraints resulting from the spatio-temporal definition of the graph. The so-called flight structure constraints define the continuity in time and space for a flight. The temporal continuity constraints are represented by the following linear inequalities and equalities:

$$x_f^{t-1} \leq x_f^t, \quad t \in T_{f,XY}, (X,Y) \in \mathcal{A}_f, f \in \mathcal{F}$$

$$x_f^t = x_f^{t-1}, \quad t \in T_{f,XY} \setminus T_{f,XY}', (X,Y) \in \mathcal{A}_f, f \in \mathcal{F}$$

(12)

where $T_{f,XY}$ is the set of feasible time units in which a flight $f$ can reach Node $Y$ from Node $X$ over the arc connecting the two nodes, and $T_{f,XY}'$ is the smallest set of consecutive time intervals that contains $T_{f,XY}$. These constraints state that if a flight was in node $X$ by time period $t$, then this must also hold true for any later time period $t' > t$.

The spatial continuity constraints are given by the following inequalities:

$$\sum_{z \in T_{f,XY}'} x_f^{t+1} \leq \sum_{x \in T_{f,XY}} x_f^{t-1}, \quad t \in T_{f,XY}, Y \in \mathcal{N}_f \setminus \{k^d, k^d\}, f \in \mathcal{F}$$

(13)

In the foregoing, $T_{f,XY}$ denotes the set of all times units by which a flight $f$ can reach Node $Y$ from any other node along the route of that flight. Spatial continuity constraints force connectivity through a node.

The third set of constraints is composed of those that are derived from airspace capacity. To formulate the problem with capacity constraints, the sets $N_{f}^+$ and $N_{f}^-$ are defined for a flight $f$ in the $j$th Sector, as the set of nodes entering and leaving the $j$th Sector. The sector capacity constraints are given by the following:

$$\sum_{j \in J} \sum_{f \in \mathcal{F}} \sum_{y \in N_{f}^{+}} \sum_{x \in T_{f,XY}} x_f^{t+1} - \sum_{y \in N_{f}^{-}} \sum_{x \in T_{f,XY}} x_f^{t-1} \leq S_{j}^{t}, \quad t \in T, j \in J$$

(14)

The foregoing equation counts the number of flights entering Sector $j$ at time $t$, and subtracts from it, the number flights leaving the Sector at that time. This number is constrained to be less than the Sector capacity at that time, $S_{j}^{t}$, for a Sector $j \in J$. Similar capacity constraints can be derived for airport arrival and departure capacity, but were not used in this work.

**C. Cost Function Formulation**

In the BLO model, the cost $J$ has contributions from different components, depending on the modeling requirements of the problem. A comprehensive list is presented in Ref. 42, which not only includes the components presented in Ref. 41, but also introduces additional terms for greater flexibility in formulating TFM problems. In this work, the number of cancelled flights, overall flight ground delays, and airborne delays were penalized. These three cost function components, denoted by $J_{\text{cancel}}, J_{\text{ground}}$, and $J_{\text{airborne}}$, are listed as follows:

$$J_{\text{cancel}} = -\sum_{f \in \mathcal{F}} \sum_{k \in \mathcal{K}} x_f^{t} I_{k,f} \quad t = \max T_{f,k} a(k_f)$$

(15)

$$J_{\text{ground}} = \sum_{f \in \mathcal{F}} \sum_{k \in \mathcal{K}} c_{f,g}(t) \left(x_f^{t} a(k_f) - x_f^{t-1} a(k_f)\right)$$

(16)
where \( c_{f,T}(t) = c_T \cdot (t - r_f) \) and \( c_{f,G}(t) = c_G \cdot (t - d_f) \) (with constant \( c_T \) and \( c_G \)) are coefficients such that each additional unit of delay from scheduled arrival and departure has a proportionately heavier penalty.

### D. Stochastic LP Formulation

The foregoing development has been used extensively for deterministic TFM. In the presence of stochastic capacity constraints, the Sector capacities \( S_f^i \) in Eq. (14) are uncertain quantities, and as described earlier can be sampled based on scenarios. The set of first stage variables for the TFM problem are then composed of all the decision variables that are unaffected by stochastic capacity. In the context of TFM, these variables correspond to the segment of a flight prior to any Sector whose capacity is stochastic. All variables that do not belong to the first stage then belong to the recourse stage. It is easy to see from Eq. (9) that the technology matrices \( T_i \) and recourse matrices \( W_i \) are respectively the same regardless of the scenario, and vectors \( h_i \) correspond to different scenarios of the capacity constraint given by Eq. (14).

### IV. A Use Case in Stochastic Traffic Flow Management

The foregoing development of the SLP is now applied to a Use Case. Figure 10 shows the area of interest, viz. the north-eastern portion of the continental United States. Five Air Route Traffic Control Centers (ARTCC) are considered in the present Use Case: Boston (ZBW), Washington, D.C. (ZDC), Indianapolis (ZID), New York (ZNY), and Cleveland (ZOB). Five airports are included in the example: Logan International Airport (BOS), Ronald Reagan Washington National Airport (DCA), Detroit Metropolitan Wayne County Airport (DTW), LaGuardia Airport (LGA), and a pseudo-airport LVT. The pseudo-airport is a waypoint from the CIFP database that lies along or near routes for flights originating from Dallas-Fort Worth International Airport (DFW) and George Bush Intercontinental Airport (IAH) that were not included in the analysis because they lie outside the geographical area of interest. The airport DTW is also located close to the jet-routes for flights originating from the west coast.

The Centers are subdivided into Sectors. This is because stochastic capacity constraints are more readily enforced at the Sector level than at the Center level, due to the availability of MAP values as a notion of capacity. The choice of Sectors in the Use Case has been restricted to those that correspond to the cruise altitude of aircraft. This results in the airspace of interest being divided into 89 Sectors, including those from Atlanta ARTCC (ZTL). Enroute Sectors from ZTL are also included in this analysis because flights from LVT to DCA, LGA, or BOS may be redirected through ZTL in the presence of inclement weather. It should be noted that the number of Sectors used in the LP formulation ultimately depends on the number of Sectors actually encountered by the modeled flights and their routes.
A. Use Case Formulation

Five weather scenarios are used to introduce stochastic Sector capacities. The weather scenarios were generated using echo tops from the NOAA website and superimposed upon the area of interest, as shown in Figure 11. The shape, size, and location of the echo tops were randomly altered in each scenario to cover different parts of the Sector layout. Each echo top is converted into a set of distinct convex polygons using k-means clustering.

Nominal trajectories are first generated, in the absence of weather. These paths are used to solve the deterministic TFM problem with rerouting and serve as the benchmark against which delays in the system are evaluated. The nominal trajectories between all airport pairs are shown in Figure 12 as blue lines. These trajectories were generated using an A* search among the waypoints from the CIPF database, using distance from the destination as the criterion. The relevant waypoints in the region of interest are also shown in this figure, as grey dots.

A composite of the five weather scenarios is depicted by the shaded area in Figure 12. It may be observed that all of the nominal trajectories intersect with Sectors that are affected by at least one weather scenario. The next step in the development of the Use Case is the identification of common paths for multiple routes for a given airport pair. Thereafter, reroutes for each weather realization are calculated from these points. The common segments correspond to the first stage of the stochastic program, and each of the reroutes starting from the end of the common segment corresponds to a set of recourse variables.

The stochastic LP is formulated for a 15 flight example, all of which depart their airports at time \( t = 1 \). The airport pairs are chosen arbitrarily and are given by: LVT-BOS, LVT-DTW, LVT-LGA, LVT-DCA, LGA-LVT, LGA-DTW, DCA-DTW, DTW-LGA, LGA-LVT, LVT-BOS, LVT-LGA, BOS-LVT, LGA-DTW, DTW-BOS, DTW-LGA. All routes between airport pairs are discretized into a network of nodes.
connected by arcs. The nodes are obtained by calculating the intersection of each route with the Sectors the flight passes through. This results in the set \( \mathcal{A}_f \) of node pairs XY.

It is assumed that all flights are cruising with a true airspeed of 400kts. Based on the length of each segment, the amount of time is calculated, i.e. the constants \( t_{f,XY} \) defined in Section IIIA. To reduce the complexity of the problem, all values are scaled by a discrete time step of 4 minutes and then rounded (floored) to the nearest integer.

The structure of the constraint matrix is shown in Figure 13. The LP is composed of 731 first stage variables and 4895 recourse variables per scenario; a total of 27208 variables (accounting for additional variables for segment-length and path-length deviations). A total of 99460 constraints are generated. The matrix is 99.97% sparse, with 337402 non-zero variables.

**Figure 13. Constraint Matrix Structure for 5 Scenarios with 15 Flights**

The matrix structure clearly exhibits the Benders’ decomposition blocks, with the first stage block (A matrix) in the top left corner, and the five T and W blocks along the rows and the diagonal. Additionally, each block exhibits the so-called Dantzig-Wolfe decomposition, with matrices along a row depicting the master problem, and fifteen blocks along the diagonal of a Bender’s block depicting the subproblem for each aircraft. Note that the subproblem blocks are of different size because the number of variables required for each flight depends on the number of segments along the path.

**B. Use Case Results**

The nominal route solver, risk-neutral case, and risk-averse cases are solved independently. Since the flight time in a link is considered fixed, i.e. \( \ell_{f,XY} = \tilde{\ell}_{f,XY} = \tilde{\ell}_{f,XY} \), path-stretching is not permitted as a solution to the LP, and no airborne delays are obtained. Instead, the solution attributes all delays to ground stops. An example of a ground stop due to downstream Sector capacity constraints is shown in Figure 14.

As mentioned before, all flights have a scheduled departure time of \( t = 1 \). In this example, two flights from LVT – one to DCA (Figure 14, right) and one to LGA (Figure 14, left) – result in a downstream Sector capacity constraint being violated. As a consequence, the flight to DCA is delayed by one time unit. This figure shows multiple trajectories; the choice of which depends on the actual scenario being realized.
Figure 14. Ground Stop at LVT

Figure 15 shows the trajectories followed by a flight from LGA to DTW, for the risk-neutral strategy (left) and the risk-averse strategy (right). In both figures, the solid blue lines depict the routes followed by a flight for all realizations of the weather, except for Scenario 3. If Scenario 3 occurs, the shaded area depicts the coverage of Sectors by inclement weather. In the risk-neutral strategy, the dashed line depicts the trajectory followed upon the realization of Scenario 3. Note that the entry and exit times over the common segment are identical. In the risk-averse case, it may be observed that flight follows the same trajectory, a route that avoids all weather-impacted areas in all scenarios. The risk-averse strategy results in the flight always requiring 22 units of time, whereas the risk-neutral trajectory can require 21 or 22 units of time, depending on which weather scenario is realized.

Figure 15. Risk-Neutral vs. Risk-Averse Strategy from LGA to DTW

C. Analysis of the Risk-Hedging Strategy

Figure 16 shows the relative frequency histogram for flight delays, for a variety of risk-hedging strategies. The delays for a strategy are calculated by subtracting the nominal flight times (obtained by solving the LP for flights with no stochastic weather constraints) from the flight times obtained from a risk-hedging LP. Over a scenario, all deviations from scheduled are summed; the result is termed the ‘system delay’. Note that the formulation used in this example penalizes all deviations from the scheduled time, including negative delays, i.e. arrivals before scheduled time. It is possible to penalize only positive delays, by removing the absolute sign in Eq. (6). In the Use Case, all delays are positive since the scheduled time of flight is equal to the time of flight along the nominal trajectory, and the time of flight on any other route other than the nominal route is larger than that for the nominal trajectory. The mean delay is calculated by adding the system delays weighted by the probability of the scenario.

In the top left, the histogram for the risk-neutral strategy is depicted. The red line in all the figures is used to depict the expected value of the delay time. The expected value of delay as a result of the risk-neutral strategy is
approximately 4 time units. The presence of an outlier results in a Maximum Absolute Deviation (M.A.D.) of approximately 20 time units.

Next, a risk-hedging constraint is imposed such that the permitted M.A.D. is 0.67 times the risk-neutral M.A.D. The resulting histogram is shown on the top right of Figure 16. It is observed that the expected delay increases to approximately 10 time units, and the M.A.D. reduces to approximately 14 time units, which is approximately 67% of 20 time units.

The bottom left of Figure 16 corresponds to a permitted M.A.D. that is 0.13 times the risk-neutral M.A.D. This strategy results in an increase in the expected value of the delay to 24 units. However, the M.A.D. reduces to 3 time units, which is approximately 13% of 20 time units.

Finally, on the bottom right, the result for the risk-averse strategy is shown. This strategy results in all flights flying along the risk-averse route that circumvents all weather phenomena. The expected delay increases to 30 time units and there is no variation in the delay since all flights follow the risk-averse trajectory regardless of which of the five weather scenarios has been realized.

![Figure 16. Relative Frequency Histogram of Deviations from Mean Delay](image)

**Figure 16. Relative Frequency Histogram of Deviations from Mean Delay**

for $\alpha = 1.00$, $\alpha = 0.67$, $\alpha = 0.13$, $\alpha = 0.00$

Although results corresponding to four values of the parameter $\alpha$ are shown in Figure 16, analysis on the Use Case was carried out by first starting with the risk-neutral value of M.A.D. and using the floor operator to set the value as risk-hedging constraint on the stochastic LP. At the end of an iteration, a new value of M.A.D. and expected delay were calculated. The parameter $\alpha$ is equal to the ratio of the new M.A.D. and the risk-neutral M.A.D. This process is iterated until the M.A.D. is equal to zero, or no feasible solution is obtained. The result of this iteration is shown in Figure 17. The solution for each iterated value of $\alpha$ is shown as a blue circle, and is a Pareto-optimal solution because it is the solution with minimum cost which satisfies the maximum absolute deviation bound. The
Pareto frontier is not necessarily convex because only integer solutions are allowed in the problem. The four strategies with different values of the iterating parameter $\alpha$ are also shown on this figure.

![Pareto Frontier and Pareto-Optimal Solutions](image)

**Figure 17. Pareto-Optimal Solutions for Trade-Off between Mean Delay and Maximum Deviation from Mean Delay**

V. Conclusions

This paper demonstrated the feasibility of a new approach to actively manage the risk induced in TFM due to the stochasticity of adverse weather. A novel approach has been developed that achieves the following: 1) a preliminary definition of risk in TFM due to weather stochasticity and demonstrated the applicability of performance-risk tradeoff concepts from Modern Portfolio Theory; 2) an approach to map the impact of stochastic adverse weather in terms of NAS capacity uncertainty; 3) risk-management strategies in the LP framework, and 4) the active control of variance in system performance, by sacrificing the expected value of cost. This guarantees the performance bounds on the TFM algorithm, while also assuring robustness in the TFM solution.

A Use Case scenario was presented that reflects NAS operations on a regional scale. The Use Case work flow can be extended to more complex networks such as a NAS-wide formulation with several thousand flights. The work advanced in this paper can be extended NAS-wide and can form the basis for the development of a decision support tool. Additional extensions include mechanisms for incorporating user preferences in the stochastic TFM. Moreover, substantial improvements in the solution speed are feasible through the use High-Performance Computing hardware. For instance, the LP formulation is amenable to decomposition techniques and parallel implementation on Graphics Processing Units. These platforms offer considerable advantages over standard LP solvers in that they can be used to solve several large scale problems with a large number of scenarios and flights in an efficient manner and in real time.

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